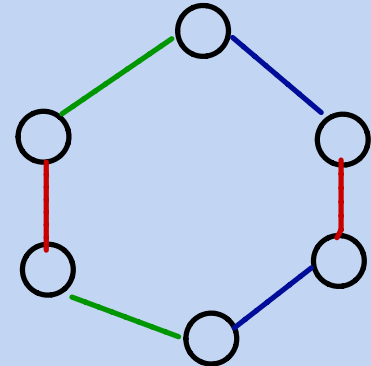
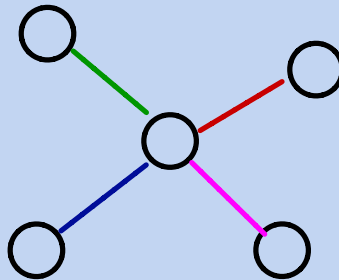
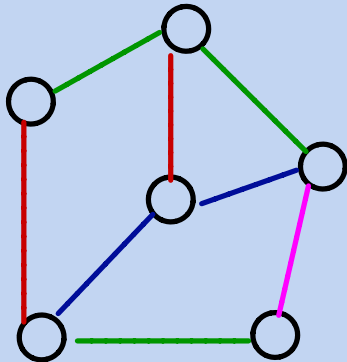


How to make a connected graph properly connected

Shinya Fujita

Yokohama City University



## Definitions:

An edge-colored graph is **properly connected**

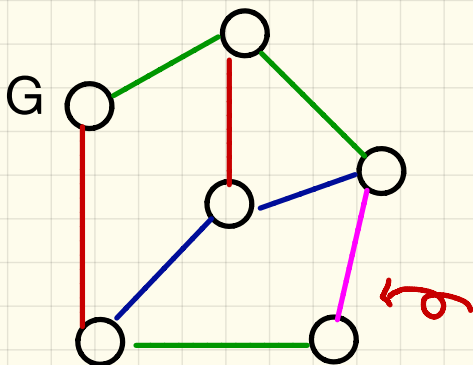
$\Leftrightarrow$  For any  $u, v \in V(G)$ ,  $\exists$  properly colored path joining them

$$\delta^c(G) := \min \{ d^c(v) \mid v \in V(G) \}$$



color degree of  $v$ ; i.e., the number of colors adjacent to  $v$  in  $G$ .

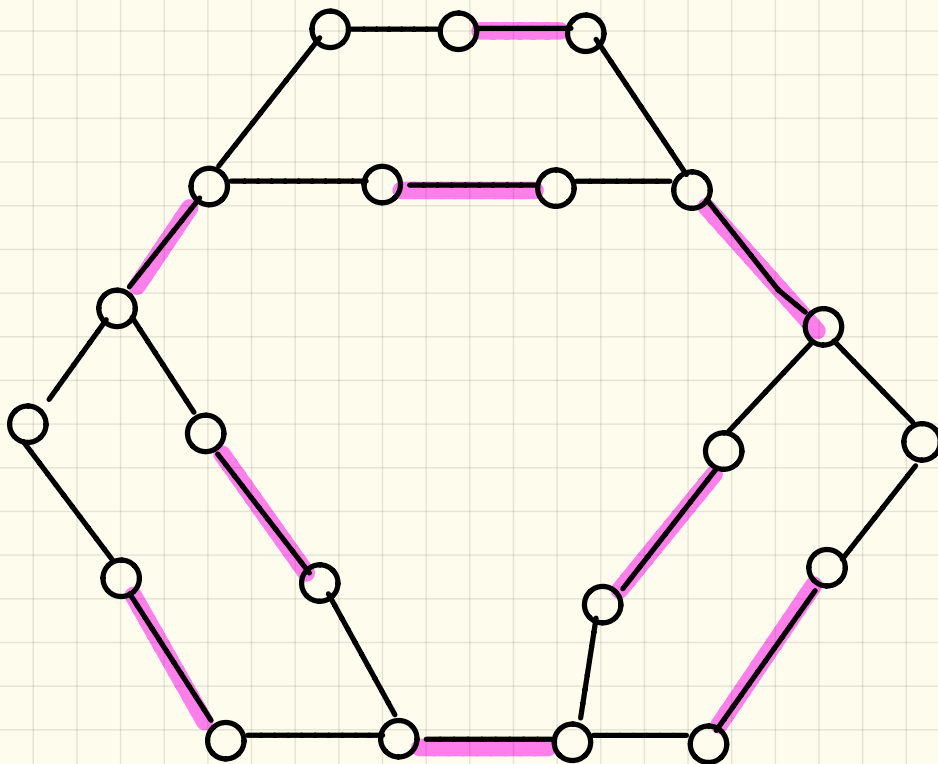
Ex.



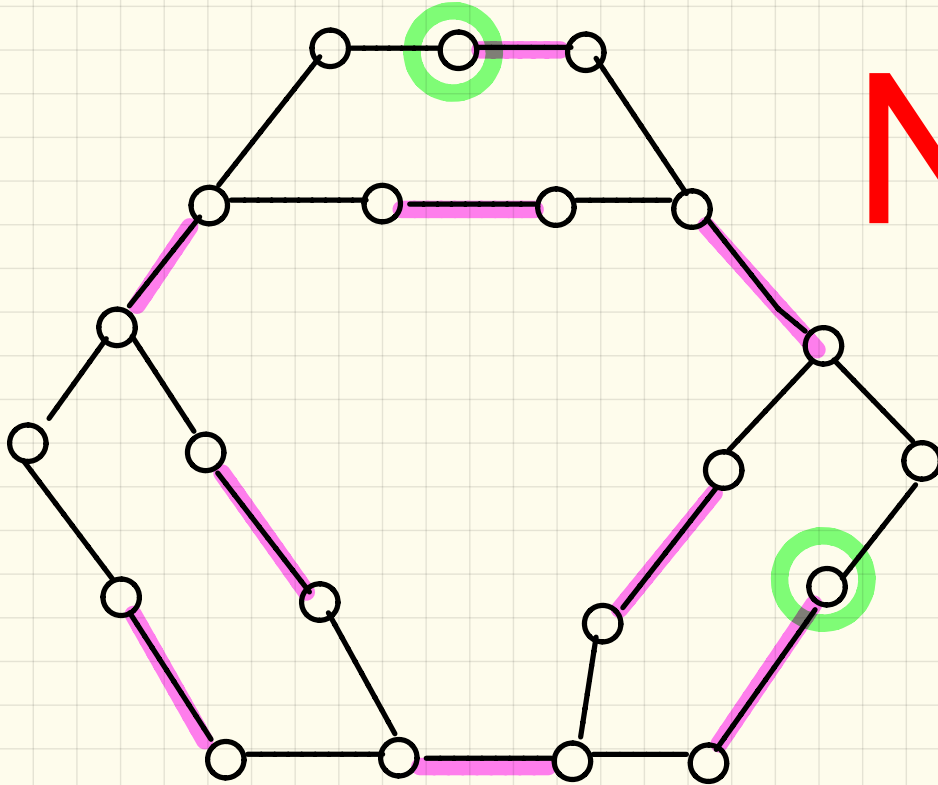
Note:  $\delta(G) \geq \delta^c(G)$

$$\delta^c(G) = 2$$

Is this properly connected ??

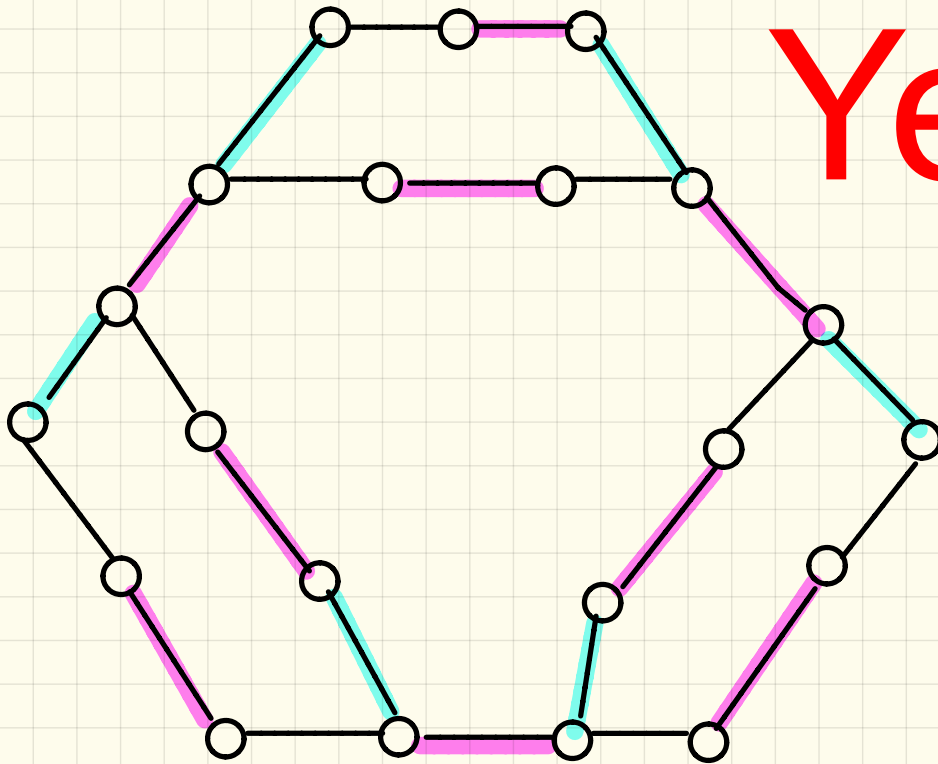


Is this properly connected ??



**No!**

Is this properly connected ??



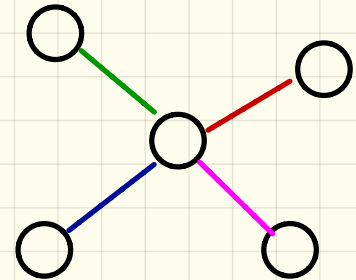
Yes!

In this talk, we consider edge-colored connected graphs.

The proper connection number of  $G$  means

$pc(G) := \min \{ c \in \mathbb{N} : \text{one can make } G \text{ properly connected by } c \text{ colors} \}$

**Note:** If  $G$  is a tree, then  
 $pc(G) \geq \Delta(G)$



Known results:

Th.1 (Magnant & F 2011)

If  $\delta'(G) \geq \frac{|V(G)|}{2}$  then  $G$  is properly connected.

Th.2 (Borožan et al. 2012)

If  $G$  is 2-connected, then  $pc(G) \leq 3$ .

Th.3 (Huang et al. 2017)

If  $G$  is 2-conn. and  $\text{diam}(G) \leq 3$ , then  $pc(G) \leq 2$ .

Th.4 (Brause et al. 2017)

If  $G$  is 2-conn. and  $\delta(G) > \max\{2, \frac{|G|+8}{20}\}$ , then  $pc(G) \leq 2$ .

Known results:

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I feel that the following is a challenging problem.

Pbm : Characterize 2-conn. graphs  $G$   
s.t.  $pc(G) = 3$ .

Th.5 (Brause et al. 2017)

For every  $d \geq 3$ , there exists a 2-conn. graph  $G$   
s.t.  $\delta(G) = d$ ,  $|G| = 42d$  and  $pc(G) = 3$ .

SPRINGER BRIEFS IN MATHEMATICS

Xueliang Li · Colton Magnant  
Zhongmei Qin

# Properly Colored Connectivity of Graphs

 Springer

In this talk, I would like to consider how to make a graph properly connected.

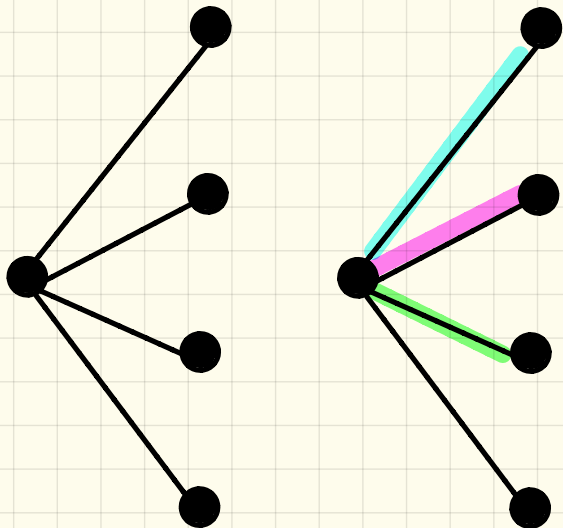
Let  $G$  be a monochromatic graph with color 1.

Define :

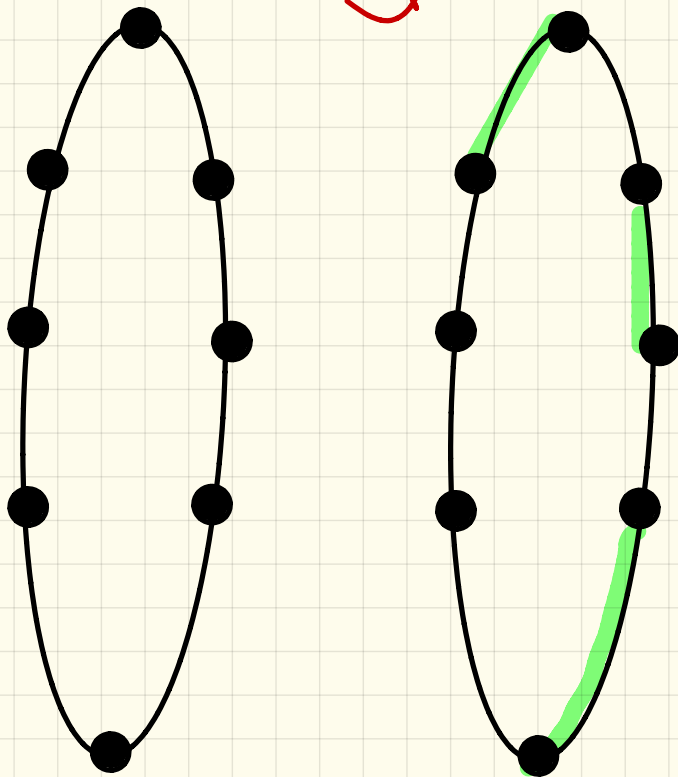
$pc_{opt}(G) = \min \{ p+q \mid \text{We can make } G \text{ properly} \}$   
connected by recoloring  
 $p$  edges with  $q$  colors.

Easy to check :

$$pc_{opt}(K_{1,m}) = 2m-2, \quad pc_{opt}(C_n) = \lfloor \frac{n-1}{2} \rfloor + 1$$

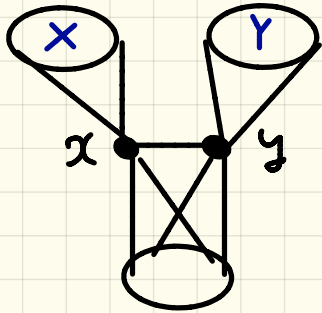


$$PC_{opt}(K_{1,4}) = 3 + 3$$



$$PC_{opt}(C_8) = 3 + 1$$

An edge  $e=xy$  is *good*  $\Leftrightarrow$

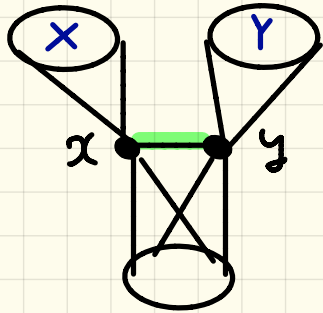


We can find a partition  
 $V(G) - (N(x) \cap N(y)) = X \cup Y$   
s.t.  $X \subset N(x)$  and  $Y \subset N(y)$ ,  
and  $G[X]$  and  $G[Y]$  are cliques.

## Th. 6 (F)

A conn. graph  $G$  has  $\chi_{\text{opt}}(G) \leq 2$  if and only if  
 $G$  contains a good edge.

An edge  $e=xy$  is *good*  $\Leftrightarrow$



We can find a partition  
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## Th. 6 (F)

A conn. graph  $G$  has  $\chi_{\text{opt}}(G) \leq 2$  if and only if  $G$  contains a good edge.

### Th.7 (F)

If  $\alpha(G) \leq 2$  then  $\rho_{\text{opt}}(G) \leq 3$ .

Th.8 (F) Let  $m \geq n \geq 2$  and  $m+n \geq 9$ .

$\rho_{\text{opt}}(K_{m,n}) = 4$  for  $n=2,3$ ; and,

$\rho_{\text{opt}}(K_{m,n}) = 5$  for  $n \geq 4$ .

### Th.9 (F)

size of a max. matching



If  $G$  is a tree, then  $\rho_{\text{opt}}(G) = |G| - 2 - \alpha'(G) + \Delta(G)$ .

### Th.10 (F)

If  $G$  is 2-conn.  $P_4$ -free with  $|G| \geq 9$ , then  $\rho_{\text{opt}}(G) \leq 5$ .

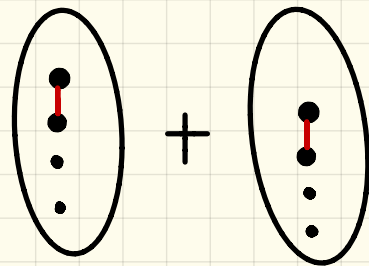
Some observations

**Prop. 1:** If  $G$  contains a complete bip. graph  $H$  s.t.

(i)  $|G|=|H|$  and (ii) each partite set contains an edge,

then  $\rho_{\text{opt}}(G) \leq 3$ .

**Cor.** If  $G$  is a complete multipartite graph  $K_{n_1, \dots, n_\ell}$   
s.t.  $\ell \geq 3$  and  $2 \leq n_1 \leq \dots \leq n_\ell$ , then  $\rho_{\text{opt}}(G) \leq 3$ .

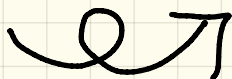
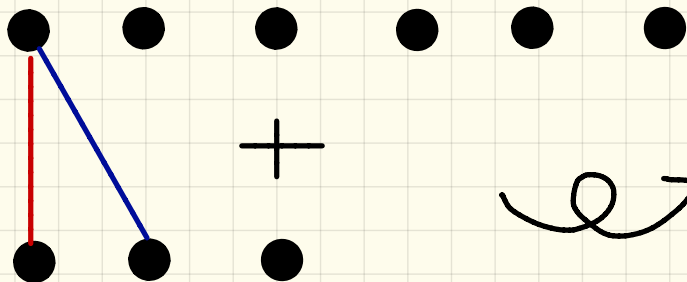




Some observations

① Upper bound on  $p_{\text{opt}}(K_{m,n})$ , where  $m \geq n \geq 2$   
and  $m+n \geq 9$

Case 1:  $2 \leq n \leq 3$

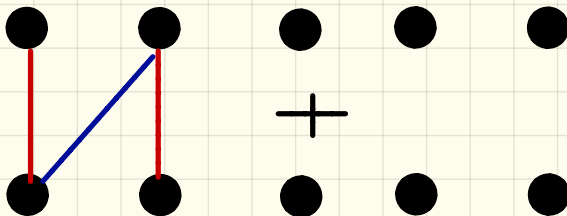
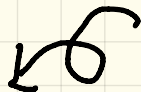


$$p_{\text{opt}}(K_{m,n}) \leq 4$$

---

Case 2:  $n \geq 4$

$$p_{\text{opt}}(K_{m,n}) \leq 5$$



## Extension

An edge colored graph  $G$  is **properly  $k$ -connected** if for any pair of vertices, there exist  $k$  internally disjoint properly connected paths joining them.

We can consider the following function of a graph.

$$\mathcal{P}_{\text{opt}}^k(G) := \min \{ p+q \mid \text{we can make } G \text{ properly } k\text{-conn.} \\ \text{by recoloring } p \text{ edges with} \\ q \text{ new colors} \}$$

## Extension

Def: A mono. conn. graph  $G$  is  $(p,q)$ -feasible

if we can make  $G$  properly connected by recoloring  
 $p$  edges with  $q$  new colors.

In particular, let's call  $(p,q)$ -optimal feasible if

$$p + q = p_{c, \text{opt}}(G).$$

Remark: If  $G$  is  $(p,q)$ -feasible, then  $p \geq q \geq p_c(G)$ .

Cor. of Th. 6:  $G$  is  $(1,1)$ -feasible  $\Leftrightarrow G$  is a complete graph

Th. 7' (F)

Any graph  $G$  s.t.  $d(G) \leq 2$  is  $(2,1)$ -feasible.

## Th. 8' (F)

For  $K_{m,n}$  with  $m \geq n \geq 2$  and  $m+n \geq 9$ ,

if  $n=2,3$ , then  $K_{m,n}$  is  $(2,2)$ -optimal feasible;  
and if  $n \geq 4$ , then  $K_{m,n}$  is  $(3,2)$ -optimal feasible.

---

## Th. 9' (F)

If  $G$  is a tree, then

$G$  is  $(n-1-\alpha'(G), \Delta(G)-1)$ -optimal feasible.